

Lecture 10

Recall $R \in \text{Rings}$, $A_R \in {}_R\text{Mod}$, $_R B \in \text{Mod}_R \rightsquigarrow A \otimes_R B \in \text{Ab}$.

Note If ${}_R A_R \in {}_R\text{Mod}_R$ (R - R bimodules),

then indeed $A \otimes_R B \in {}_R\text{Mod}$.

→ can check details on your own

e.g. $V, W \in \text{Vect}_k \Rightarrow V \otimes W \in \text{Vect}_k$ as well

Maps

prop. let $f: A_R \rightarrow A'_R$ (morphism in Mod_R)

$g: {}_R B \rightarrow {}_{R'} B'$ ($___ \text{ in } {}_R \text{Mod}$)

Then $\exists!$ \mathbb{Z} -map denoted $f \otimes g: A \otimes_R B \rightarrow A' \otimes_{R'} B'$

such that $f \otimes g: a \otimes b \mapsto f(a) \otimes g(b)$.

Recall the pure tensors $\{a \otimes b \mid a \in A, b \in B\}$ generate $A \otimes_R B$.

Pf.

$$\begin{array}{ccc} A \times B & \xrightarrow{h} & A \otimes B \\ (a, b) \swarrow \varphi \quad \searrow \exists! \tilde{\varphi} & & \text{call this map "f} \otimes g\text{"} \\ & A' \otimes B' & \\ & f(a) \otimes g(b) & \end{array}$$

STS φ is R -biadditive.

- additivity on each side is clear ✓

- $(ar, b) \xrightarrow{\varphi} f(ar) \otimes g(b)$

$$= f(a)r \otimes g(b)$$

$$= f(a) \otimes rg(b)$$

$$= f(a) \otimes g(rb) \xleftarrow{\varphi} (a, rb)$$

b/c f, g are R -linear.

□

We can also compose these maps b/w tensor products:

Cor. Given $A \xrightarrow{f} A' \xrightarrow{f'} A''$ (in Mod_R)
 $B \xrightarrow{g} B' \xrightarrow{g'} B''$ (in $R\text{Mod}$)

we have $(f' \otimes g') \circ (f \otimes g) = f'f \otimes g'g$

Pf. Because of the uniqueness of the map lifting
 $\varphi: a \otimes b \mapsto f'f(a) \otimes g'g(b)$. □

(Recall)

defn. Let \mathcal{C}, \mathcal{D} be categories.

A functor $F: \mathcal{C} \rightarrow \mathcal{D}$ is a function such that

- if $A \in \mathcal{C}$, then $F(A) \in \mathcal{D}$.
- If $f \in \underline{\mathcal{C}(A, A')}$, then $F(f) \in \underline{\mathcal{D}(F(A), F(A'))}$

Common

notation for

$$\underline{\mathcal{M}\>_{\mathcal{C}}(A, A')} = \underline{\text{Hom}_{\mathcal{C}}(A, A')}$$

- F preserves composition of morphisms

$$F(gf) = F(g)F(f)$$

- F preserves identity morphisms

$$F(\text{id}_A) = \text{id}_{F(A)}$$

Thm. Fix $A_R \in \text{Mod}_R$.

There is an additive functor

$$(A \otimes -) : {}_R\text{Mod} \rightarrow \text{Ab}$$

$$(\text{Ob}) \quad {}_R B \mapsto A \otimes_R B$$

$$(\text{Mor}) \quad [g: B \rightarrow B'] \mapsto [\text{id}_A \otimes g: A \otimes B \rightarrow A \otimes B']$$

covariant
functor!

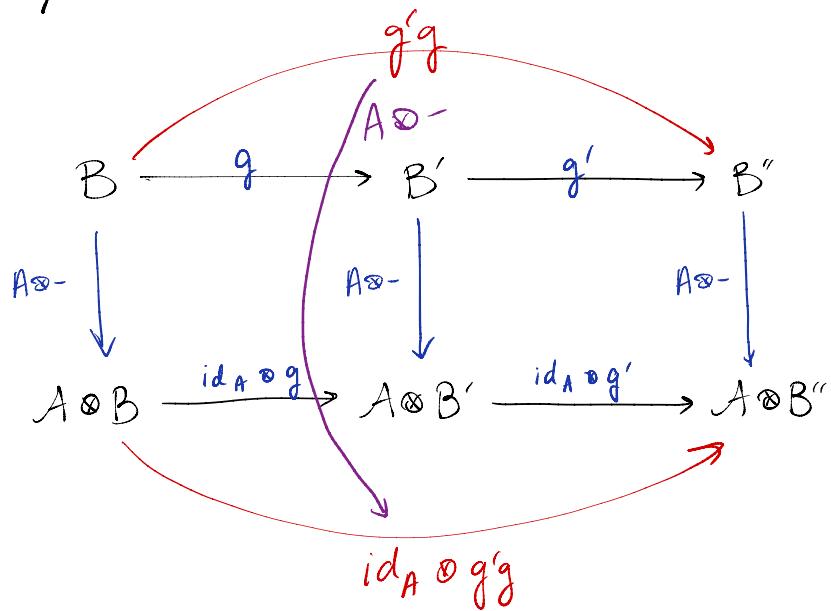
Pf.

① $A \otimes -$ is a functor:

$$\text{identities are preserved: } \text{id}_A \otimes \text{id}_B = \text{id}_{A \otimes B}$$

(by uniqueness in CP maps)

composition:



ok, by above corollary ✓

② Additive functor:

Abelian cat: hom sets are abelian groups

$$F = A \otimes -$$

NTS $F(g+h) = F(g) + F(h)$:

$$\text{id}_A \otimes (g+h) = \text{id}_A \otimes g + \text{id}_A \otimes h$$

✓ since both sides send the generators

$$a \otimes b \mapsto a \otimes g(b) + a \otimes h(b).$$



Rmk. Same construction for $- \otimes_R B : \text{Mod}_R \rightarrow \text{Ab}$.

Cov. If $f: M \xrightarrow{\cong} M'$ and $g: N \xrightarrow{\cong} N'$ are isoms.

then so is $f \otimes g: M \otimes N \rightarrow M' \otimes N'$.

Pf.

$$f \otimes g = (f \otimes \text{id}_{N'}) (\text{id}_M \otimes g).$$

+

Rmk.

$$\textcircled{1} \quad M \in {}_R\text{Mod} \Rightarrow R \otimes_R M \cong M$$

$r \otimes m \mapsto rm$

$$\textcircled{2} \quad M \in \text{Mod}_R \Rightarrow M \otimes R \cong M$$

$$\textcircled{3} \quad A \otimes \left(\bigoplus_{i \in I} B_i \right) \cong \bigoplus_{i \in I} A \otimes B_i$$

(Again, uniqueness of the map in the UP)

$$A \times \left(\bigoplus_i B_i \right) \longrightarrow A \otimes \left(\bigoplus_i B_i \right)$$

(a, b_1, b_2)

\downarrow

$\exists!$

$a \otimes (b_1 + b_2) = a \otimes b_1 + a \otimes b_2$

$\bigoplus_i A \times B_i$

$a \otimes b_1 + a \otimes b_2$

Note: finite sum everywhere

(4) If k is commutative, then $A \otimes_k B \cong B \otimes_k A$

Key: May regard $M \otimes M$ as a $(k-k)$ -bimodule by enforcing $aM = ma$.

$$\text{Then } A \times B \longrightarrow A \otimes B \quad a \otimes b$$

Next time: Back to the story of Proj. injective, ...
etc.

thm. $A \otimes -$ is a (covariant) right-exact functor.